

J-Logic: Logical Foundations for JSON Querying

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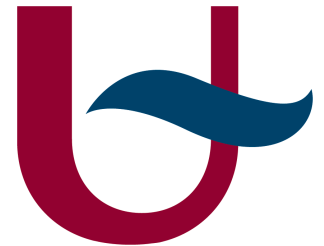
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Logical Foundations for JSON Querying

- **Object descriptions:** sets of path–value pairs
- **Path variables**
- **Generating new keys** by packing
- Testing the **object-to-object property**
- **Containment** testing
- Open problems
- JSON-to-JSON queries,
easy access to deeply nested objects

Describing a JSON object...

```
{ name:john,  
  children: { 1:{name:anne,age:12},  
              2:{name:bob,age:18},  
              3:{name:chris,age:24} } }
```

...as a set of path–value pairs:

```
{ name:john,  
  children.1.name:anne,  
  children.1.age:12,  
  children.2.name:bob,  
  children.2.age:18,  
  children.3.name:chris,  
  children.3.age:24 }
```

“object description”

J-Logic ingredients

- Object descriptions as binary relations
- Input as well as output
- Path variables $\$x$, concatenation .
- Atomic variables $@x$
- Atomic value constants $john$, 42
- Logic-based language (Datalog)

[Sequence Datalog, Bonner&Mecca]

Hello World in J-Logic

T(hello:world).

> T={hello:world}

My first real J-Logic program

```
> Sales={ 1:{1999:40,2017:35},
          2:{1999:70,2001:80},
          3:{2001:90,2002:70,2017:33}}
```

```
ByYear(@y.@p:@s) :- Sales(@p.@y:@s).
```

```
> ByYear = { 1999:{1:40,2:70},
             2001:{2:80,3:90},
             2002:{3:70},
             2017:{1:35,3:33}}
```

> Sales={ 1:{1999:40,2017:35},
 2:{1999:70,2001:80},
 3:{2001:90,2002:70,2017:33}}



> ByYear = { 1999:{1:40,2:70},
 2001:{2:80,3:90},
 2002:{3:70},
 2017:{1:35,3:33}}

object descriptions!

ByYear(@y.@p:@s) :- Sales(@p.@y:@s).

ByYear={1999.1:40,	Sales={1.1999:40,
2017.1:35,	1.2017:35,
1999.2:70,	2.1999:70,
2001.2:80,	2.2001:80,
2001.3:90,	3.2001:90,
2002.3:70,	3.2002:70,
2017.3:33}	3.2017:33}

Deep equality in J-Logic

> R={ a: ,
b:  }

// Are o_1 and o_2 equal?

T(answer:no) :- R(a.\$x:@v), not R(b.\$x:@v).

T(answer:no) :- R(b.\$x:@v), not R(a.\$x:@v).

Q(answer:no) :- T(answer:no).

Q(answer:yes) :- not T(answer:no).

Key generation by packing

```
// input: R, a deeply nested object  
// retrieve all subobjects with loc:chicago  
Q(<$x>.$z:@v) :- R($x.loc:chicago),  
                R($x.$z:@v).
```

```

// input: R, a deeply nested object
// retrieve all subobjects with loc:chicago
Q(<$x>.$z:@v) :- R($x.loc:chicago),
                R($x.$z:@v).

> S={a:1,b:2}
> T={a:3,b:4}
// cartesian product of JSON objects
Cart(<@x.@y>.s.@x:@u) :- S(@x:@u),T(@y:@v).
Cart(<@x.@y>.t.@y:@v) :- S(@x:@u),T(@y:@v).
> Cart={<a.a>:{s:{a:1},t:{a:3}},
        <a.b>:{s:{a:1},t:{b:4}},
        <b.a>:{s:{b:2},t:{a:3}},
        <b.b>:{s:{b:2},t:{b:4}}}}

```

Cartesian product by packing

Proper objects

- JSON objects are unordered
- Are `{loc:raleigh,loc:chicago}` and `{loc:chicago,loc:raleigh}` the same?
- Not in JavaScript!
- Objects where keys are not keys (!) are called **improper**
- Input objects are normally proper
- We should avoid outputting improper objects
- **Object-object property: proper inputs** are mapped to **proper outputs**

The object-object (O2O) property

> $R = \{\text{raleigh}:\text{loc}, \text{chicago}:\text{loc}\}$

$Q(@v:@k) \text{ :- } R(@k:@v)$. // not O2O

> $Q = \{\text{loc}:\text{raleigh}, \text{loc}:\text{chicago}\}$

- All other example programs we have seen so far are O2O
- **Theorem:** It is decidable (in EXPTIME) whether a given **positive, recursion-free** J-Logic program is O2O
- “Unions of conjunctive queries” for J-Logic

Deciding the O2O property

- “JSON atomic equality-generating dependencies” (**jaegd**)
- D is proper if it satisfies these two jaegds:
 $D(\$x:@u), D(\$x:@v) \rightarrow @u=@v$
 $D(\$x:@u), D(\$x.\$y:@v) \rightarrow \text{false}$
- Plan:
 1. Reduce O2O decision to the **implication problem for jaegds**
 2. Solve the jaegd implication problem

Chasing jaegds?

- Complications due to two kinds of variables
- Consider Σ :

$R(@x:42) \rightarrow \text{false}$

$R(<\$x>:42) \rightarrow \text{false}$

$R(\$x.\$y:42) \rightarrow \text{false}$

- Then Σ logically implies σ :

$R(\$x:42) \rightarrow \text{false}$

- However, chasing σ with Σ does not fail
- Fortunately, dependencies expressing properness do not have this problem

Unification in J-Logic

E.g. $\$x.\$y = \$z.\w

Split rule 3 ways:

$\$x$	$\$y$
$\$z$	$\$w$

$\$x/\$z, \$y/\w

$\$x$	$\$y$
$\$z$	$\$w$

$\$x.\$s/\$z, \$s.\$w/\y

$\$x$	$\$y$
$\$z$	$\$w$

$\$z.\$s/\$x, \$s.\$y/\w

The containment problem for J-Logic

- **Theorem:** Containment of a conjunctive J-Logic query in a union of conjunctive J-Logic queries, **over flat instances**, is decidable (Π -P-2)
- Flat instances have no packed keys
- We solve the **inclusion problem** for **pattern languages** (over an infinite alphabet) extended with atomic variables.

Conclusions

- **Path variables & packing**, useful for declarative JSON querying
- Further research:
 - Query processing
 - Theory of egds–tgds for J-Logic
 - Precise complexity of O2O problem, containment
 - Containment over non-flat instances?
 - Is packing necessary for recursion-free **flat-flat** queries?



Pattern inclusion

- $\$x.\y included in $@x.\$y$
- No homomorphism from right to left!
- Replace left pattern by four variants:
 - $@x_1.@y_1$
 - $@x_1.@x_2.@y_1$
 - $@x_1.@y_1.@y_2$
 - $@x_1.@x_2.@y_1.@y_2$