



Query languages for matrices and K -relations

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Relational databases

$F =$

<i>name</i>	<i>friend</i>
john	william
john	thomas
william	francis
	...

$B =$

<i>name</i>	<i>year</i>
john	1967
thomas	1965
francis	1974
william	1980
	...

Database instance: relational structure

- assign domains to attributes
- compatible attributes have same domain
- assign sets of tuples to relation names

Database schema: $F(name, friend), B(name, year)$

Relational algebra

Union \cup

Difference $-$

Selection $\sigma_{P(A_1, \dots, A_k)}$

e.g. $\sigma_{year \geq 2000}(B)$

Natural join \bowtie

Generalized projection $\pi_f(A_1, \dots, A_k)$

e.g. $\pi_{year-2000}(B)$

Renaming $\rho_{A/B}$

Matrix databases

Data science

$$A = \begin{pmatrix} 5 & 2 & 0 \\ 2 & 1 & 3 \end{pmatrix} \qquad B = \begin{pmatrix} 300 \\ 250 \\ 330 \end{pmatrix}$$

Matrix schema uses size symbols: $A(\alpha \times \beta)$, $B(\beta \times 1)$

MATLANG

1

$$1 \begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}$$

MATLANG

1

diag

$$\text{diag} \begin{pmatrix} 1 \\ 2 \\ 3 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 2 & 0 \\ 0 & 0 & 3 \end{pmatrix}$$

MATLANG

1

diag

(conjugate) transpose

matrix multiplication

pointwise functions $f(M_1, \dots, M_k)$

e.g. $\begin{pmatrix} 1 & 2 \\ 3 & 4 \end{pmatrix} \circ \begin{pmatrix} 4 & 3 \\ 2 & 1 \end{pmatrix} = \begin{pmatrix} 4 & 6 \\ 6 & 4 \end{pmatrix}$ pointwise multiplication

Some simple MATLAB tricks

Let A be the adjacency matrix of a graph on n nodes

Number of nodes:

$$N = \mathbf{1}(A)^* \cdot \mathbf{1}(A)$$

Degree vector, duplicated n times:

$$D = A \cdot \mathbf{1}(A) \cdot \mathbf{1}(A)^*$$

Google matrix:

$$G_{ij} = d \frac{A}{D} + \frac{1-d}{N}$$

Our proposal

$$\frac{\text{relational algebra}}{\text{relational databases}} = \frac{\text{MATLANG}}{\text{matrix databases}}$$

- What is the precise expressive power?
- How does it compare to relational database querying?

Matrix database as relational database

$$A = \begin{pmatrix} 7 & 8 & 9 \\ 10 & 11 & 12 \end{pmatrix}$$

$$B = \begin{pmatrix} 300 \\ 250 \\ 330 \end{pmatrix}$$

$A =$

<i>row</i>	<i>col</i>	<i>K</i>
1	1	7
1	2	8
1	3	9
2	1	10
2	2	11
2	3	12

$B =$

<i>row</i>	<i>K</i>
1	300
2	250
3	330

K -relations: generalization of the relational database model

Every tuple is annotated with a value from some fixed semiring K

(Positive) relational algebra on K -relations

Influential paper from 2007 [Green, Garvounarakis, Tannen]

Union: adds annotations

Natural join: multiplies annotations

Selection $\sigma_{A=B}$: sets annotations to 0 for non-qualifying tuples

Projection π_{A_1, \dots, A_k} : sums annotations

Renaming

1: sets annotations to 1

Theorem

ARA(3): Annotation-Relation Algebra, width ≤ 3

Assume K is commutative

Matrix query, expressible in ARA(3) if and only expressible in MATLANG with only $+$ and \circ as pointwise functions

Analogue to classical result by Tarski and Givant:

Our result	Tarski and Givant
matrix queries	binary-relation queries
ARA(3)	FO(3)
MATLANG	classical algebra of binary relations

Expressiveness limitations of MATLANG

Not expressible in MATLANG:

- transitive closure of a graph
- testing for 4-clique

Adding matrix inverse to MATLAB

Akin to solving a system of linear equations

Expressible in MATLAB + inverse:

- PageRank vector of a graph:

$$\frac{1-d}{n} \left(I - d \frac{A}{D} \right)^{-1} \cdot \mathbf{1}$$

(by definition)

- transitive closure: let $B = A/(n+1)$

$$\sum_{k=0}^{\infty} B^k = (I - B)^{-1}$$

- number of connected components, testing bipartiteness

Eigenvectors

Eigen-decomposition, another workhorse in data analysis

Diagonalizable $A = B \cdot \Lambda \cdot B^{-1}$ where B is a basis of eigenvectors of A

Λ has the eigenvalues on the diagonal

Define: $\text{eigen}(A) := B$, nondeterministic!

Theorem: Inverse is expressible in $\text{MATLANG} + \text{eigen}$

Open problem: Show a graph query that is:

- deterministically expressible in $\text{MATLANG} + \text{eigen}$
- not in $\text{MATLANG} + \text{inverse}$

References

Brijder, Geerts, Van den Bussche, Weerwag *On the expressive power of query languages for matrices*, ICDT 2018

- Full version in TODS
- Research highlight, SIGMOD Record

Brijder, Gyssens, Van den Bussche *On matrices and K -relations*, to appear

Floris Geerts *On the expressive power of linear algebra on graphs*, ICDT 2019

Related work: LaraDB, SQL, in-database machine learning, etc.